



To each arc  $(i, j)$  of the graph  $G = (N, A)$  we associate a variable  $\tau_{ij}$  called artificial pheromone trail, shortened to pheromone trail in the following. Pheromone trails are read and written by the ants. The amount (intensity) of a pheromone trail is proportional to the utility, as estimated by the ants, of using that arc to build good solutions.

Each ant builds, starting from the source node, a solution to the problem by applying a step-by-step decision policy (Eq. 1). At each node, local information stored on the node itself or on its outgoing arcs is read (sensed) by the ant and used in a stochastic way to decide which node to move to next. At the beginning of the search process, a constant amount of pheromone (e.g.,  $\tau_{ij} = 1, \forall (i, j) \in A$ ) is assigned to all the arcs. When located at a node  $i$  an ant  $k$  uses the pheromone trails  $\tau_{ij}$  to compute the probability of choosing  $j$  as next node:

$$p_{ij}^k = \begin{cases} \frac{\tau_{ij}^\alpha}{\sum_{l \in N_i^k} \tau_{il}^\alpha}, & \text{if } j \in N_i^k; \\ 0, & \text{if } j \notin N_i^k; \end{cases} \quad (1)$$

where  $N_i^k$  is the neighborhood of ant  $k$  when in node  $i$ . Where  $\alpha$  is the parameter which determines the relative influence of the pheromone trail and the heuristic information. In S-ACO the neighborhood of a node  $i$  contains all the nodes directly connected to node  $i$  in the graph  $G = (N, A)$ , except for the predecessor of node  $i$  (i.e., the last node the ant visited before moving to  $i$ ). In this way the ants avoid returning to the same node they visited immediately before node  $i$ . Only in case  $N_i^k$  is empty, which corresponds to a dead end in the graph, node  $i$ 's predecessor is included into  $N_i^k$ .

#### Ant System (AS)

The two main phases of the AS algorithm constitute the ants' solution construction and the pheromone update. In AS a good heuristic to initialize the pheromone trails is to set them to a value slightly higher than the expected amount of pheromone deposited by the ants in one iteration.

In AS,  $m$  (artificial) ants are put on randomly chosen nodes. At each construction step, ant  $k$  applies a probabilistic action choice rule, called random proportional rule, to decide which city to visit next. In particular, the probability of ant  $k$ , currently at node  $i$ , chooses to go to node  $j$  is

$$p_{ij}^k = \frac{[\tau_{ij}]^\alpha [\eta_{ij}]^\beta}{\sum_{l \in N_i^k} [\tau_{il}]^\alpha [\eta_{il}]^\beta}, \text{ if } j \in N_i^k, \quad (2)$$

Where  $\eta_{ij} = 1/d_{ij}$  is a heuristic value that is available a priori,  $d$  is the distance between nodes,  $\alpha$  and  $\beta$  are two parameters which determine the relative influence of the pheromone trail and the heuristic information, and  $N_i^k$  is the feasible neighborhood of ant  $k$  when being at city  $i$ , that is, the set of cities that ant  $k$  has not visited yet. By this probabilistic rule, the probability of choosing a particular arc  $(i, j)$  increases with the value of the associated pheromone trail  $\tau_{ij}$  and of the heuristic information value  $\eta_{ij}$ . The role of the parameters  $\alpha$

and  $\beta$  is the following. If  $\alpha = 0$ , the closest cities are more likely to be selected: this corresponds to a classic stochastic greedy algorithm (with multiple starting points since ants are initially randomly distributed over the cities). If  $\beta = 0$ , only pheromone amplification is at work, that is, only pheromone is used, without any heuristic bias. This generally leads to rather poor results and, in particular, for values of  $\alpha > 1$  it leads to the rapid emergence of a stagnation situation, that is, a situation in which all the ants follow the same path and construct the same tour, which, in general, is strongly suboptimal.

Each ant  $k$  maintains a memory  $M^k$  which contains the cities already visited, in the order they have visited. This memory is used to define the feasible neighborhood  $N_i^k$  in the construction rule given by equation (2). In addition, the memory  $M^k$  allows ant  $k$  both to compute the length of the tour  $T^k$  it generated and to retrace the path to deposit pheromone.

After all the ants have constructed their tours, the pheromone trails are updated. This is done by first lowering the pheromone value on all arcs by a constant factor, and then adding pheromone on the arcs the ants have crossed in their tours. Pheromone evaporation is implemented by

$$\tau_{ij} \leftarrow (1 - \rho)\tau_{ij}, \quad \forall (i, j) \in L \quad (3)$$

Where  $0 < \rho \leq 1$  is the pheromone evaporation rate. The parameter  $\rho$  is used to avoid unlimited accumulation of the pheromone trails and it enables the algorithm to "forget" bad decisions previously taken. In fact, if an arc is not chosen by the ants, its associated pheromone value decreases exponentially in the number of iterations. After evaporation, all ants deposit pheromone on the arcs they have crossed in their tour:

$$\tau_{ij} \leftarrow \tau_{ij} + \sum_{k=1}^m \Delta \tau_{ij}^k, \quad \forall (i, j) \in L \quad (4)$$

Where  $\Delta \tau_{ij}^k$  is the amount of pheromone ant  $k$  deposits on the arcs it has visited.

It is defined as follows:

$$\Delta \tau_{ij}^k = \begin{cases} 1/C^k, & \text{if arc } (i, j) \text{ belongs to } T^k; \\ 0, & \text{otherwise;} \end{cases} \quad (5)$$

Where  $C^k$ , the length of the tour  $T^k$  built by the  $k^{\text{th}}$  ant, is computed as the sum of the lengths of the arcs belonging to  $T^k$ . By means of equation (5), the better an ant's tour is, the more pheromone the arcs belonging to this tour receive. In general, arcs that are used by many ants and which are part of short tours, receive more pheromone and are therefore more likely to be chosen by ants in future iterations of the algorithm.

### III. PROBLEM DESCRIPTION

In this paper a 2-D square map overlaid with a uniform pattern of grid points has been considered. The size of a map can be changed arbitrarily; here, the map consists of a 20 x 20 grid. The left bottom corner of the map is the starting point for a path while the right top corner of the map is the destination point for a path. The shape of an obstacle is always a circle, but the size of the obstacle is variable from 1 to 5 grid points. The positions of the obstacles are randomly selected and can be located at any grid point in the map except at points close to the starting point region or close to the goal point region. Furthermore, multiple obstacles are possible. FIGURE 2 shows an example of such an arrangement.

### IV. PROPOSED SOLUTION

The Travel salesman problem (TSP) is a paradigmatic optimization problem since it is used to demonstrate the original AS (Ant System) problem. Since then it has often been used as a benchmark to test the new ACO concept. The proposed path-finding algorithm employs the original concept of the ACO algorithm with some modification. The idea of the proposed algorithm is as follows.

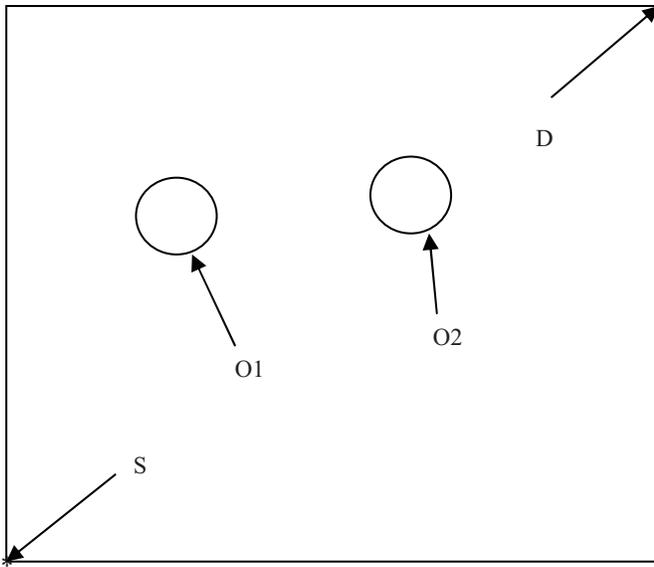


FIGURE 2: Problem configuration (S: starting point, O1, O2: obstacles, D: destination point)

Starting from the grid point (0, 0), an ant iteratively moves from a grid point to one of its neighboring grid points. When at the  $i$ th grid point (x, y), ant k can choose the next (jth) grid point by choosing one of its 8 neighbor locations: [(x+1, y+1), (x+1, y), (x, y-1), .....]  $\in N$ , where N in general is the set of all neighboring locations of the current location.

The ant takes its next step randomly, based on the probability given by

$$P_{ij}^k(t) = \frac{[\tau_{ij}]^\alpha [\eta_{ij}]^\beta}{\sum_{l \in N_i^k} [\tau_{il}]^\alpha [\eta_{il}]^\beta} + \alpha * \Omega_{ij} \quad (6)$$

Where  $\tau_{ij}(t)$  is accumulated pheromone on the jth grid point when the  $i$ th grid point is the ant's current location at time t. The quantity  $\tau_{il}$  indicates every possible lth neighbor point when the ant is in  $i$ th position.  $\Omega_{ij}$  is the dot product of the vector from  $i$  to  $j$  with the vector  $i$  to destination point. The first term is associated with the pheromone amount in the 8 neighboring grid points (a local pheromone). The second term is associated with a global attraction, where  $\alpha$  is a scale parameter. The global term is defined as the dot product of the agent's heading direction and the food (destination) direction. The purpose of the global term is to guide the agent in the desired direction so that a large number of ants can reach the goal point in a limited number of steps.

TABLE 1: An agent's current position (bold) with 8 possible next positions

(x-1, y+1)	(x, y+1)	(x+1, y+1)
(x-1, y)	<b>(x, y)</b>	(x+1, y)
(x-1, y-1)	(x, y-1)	(x+1, y-1)

The solution construction ends after each ant reaches the destination or the ant took too many steps (50 or 100 steps, for example). However, the pheromone will be deposited only if the ant reaches to the destination position in less than a certain number of steps. Thus, after a group of agents finished its tour, the pheromone in the entire map will be updated by

$$\tau_{ij}(t+1) = (1 - \rho) * \tau_{ij}(t) + \sum_{k=1}^m \Delta\gamma_{ij}^k(t) \quad (7)$$

Where  $0 < \rho < 1$  is the pheromone forgetting parameter. This parameter prevents the map from unlimited accumulation of the pheromone. The quantity  $\Delta\gamma_{ij}^k$  is the amount of pheromone that ant k deposits on the map. It is defined as

$$\Delta\gamma_{ij}^k(t) = \frac{1}{L^k(t)} \quad (8)$$

Where  $L^k(t)$  is the length of the kth ant's tour (path). When the kth ant cannot reach the destination in a certain number of steps, let the  $L^k(t) = \text{Infinity}$ .

### V. RESULTS DISCUSSION

The experiments are done with different number of obstacles, different population sizes, and a different number of iterations. Furthermore, these studies include the effects of elitism. The circular objects represent the obstacles. The obstacles can be located at any place on the map except at the starting point and at the destination point. The experiments are conducted using MATLAB.

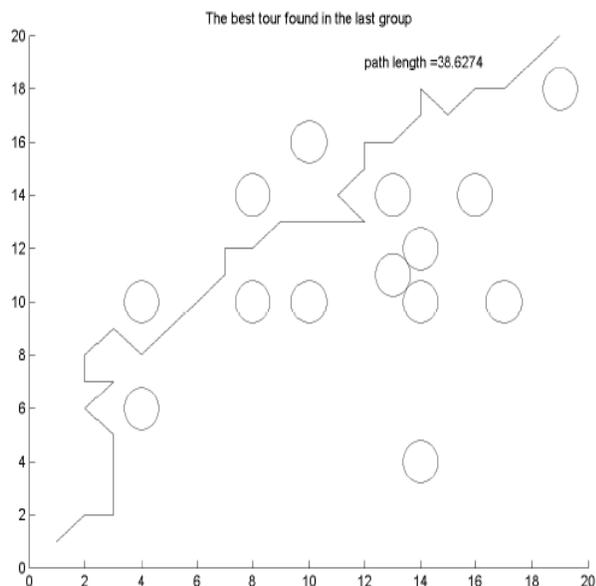


FIGURE 3: Path followed by 100 robots avoiding multiple obstacles

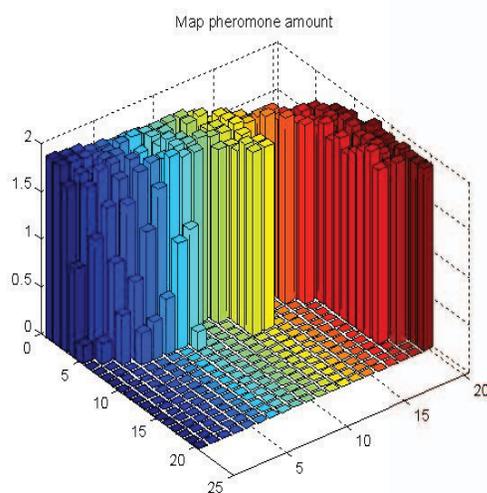


FIGURE 4: Pheromone accumulation in 3D

The simulated robot path planning algorithm generated and set the obstacles randomly. FIGURE 3 shows the result obtained from simulation, (0, 0) is start node and (20, 20) is target node. The optimum path is shown with line. From FIGURE 3 it is clear that with the help of proposed approach robots are able to find the optimal path by avoiding the obstacles.

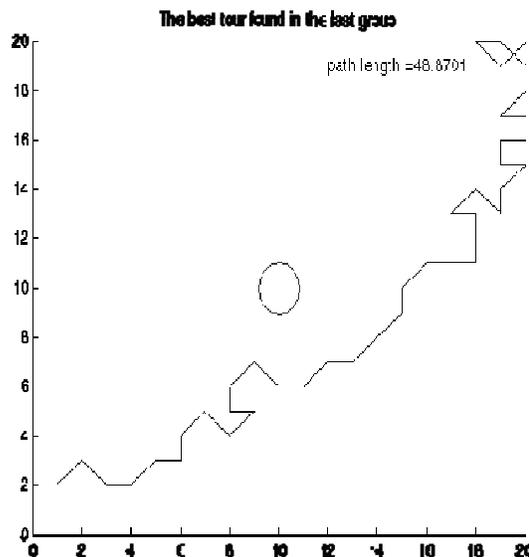
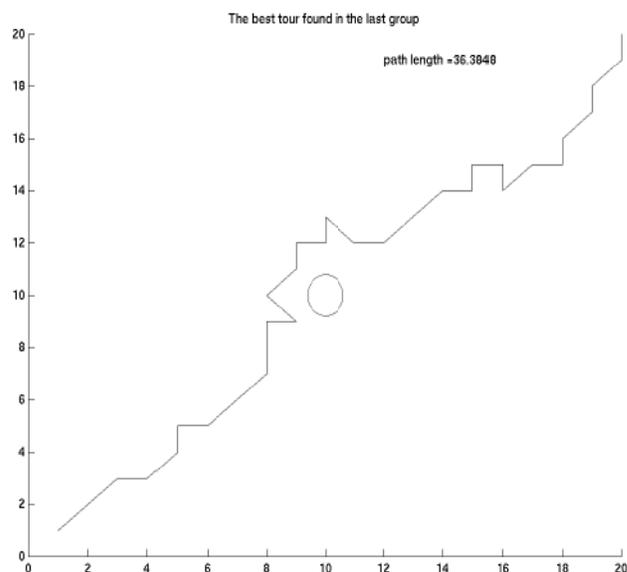


FIGURE 5: Path followed by 10 numbers of ants

FIGURE 6: Path followed by 100 numbers of ants



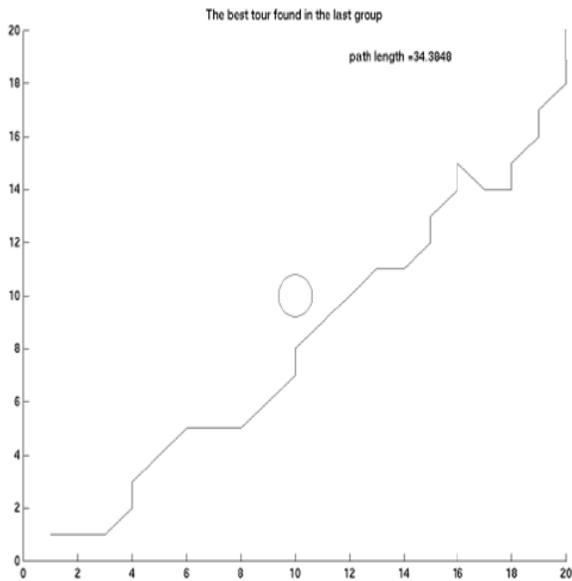


FIGURE 7: Path followed by 500 numbers of ants

TABLE 2: Effect of number of agents

The number of ants	10	100	500
Average path length of 50 trials	48.8701	36.3848	34.3848

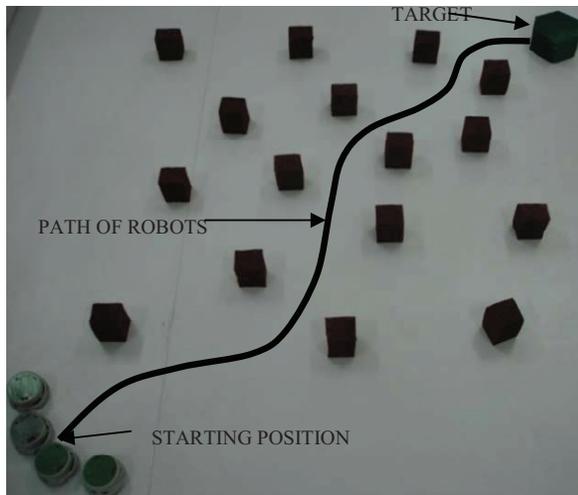


FIGURE 8: Path followed by Mobile robots during Experiment

From the results (FIGURE 5, FIGURE 6 and FIGURE 7) it is observed that a larger size group is better than a smaller size group. With the increase of numbers of Robots, a better optimized path is obtained (resulting less distance travel to reach the goal). If the group size is too small, not only does the number of ants that can reach the destination get smaller but also the amount of pheromone accumulation decreases. FIGURE 8 shows the Path of four Khepera II Mobile Robots during experiment. The Simulation results and Experimental results are in good agreement.

## VI. CONCLUSION

From the above results and discussion it is concluded that with the help of developed ACO algorithm robots are able to find the target in optimal path by avoiding the obstacles. During comparison of the results it is concluded that a larger size group of robots is better than a smaller size group of robots. In this paper, an algorithm based on ant colony is proposed to solve robot path planning problems. The proposed ant algorithm has the following characteristics: (1) it hurdles the abuse of local optimization and expedite searching speed, with the experiments, the method is proved effective. (2) The algorithm has good scalability. This is because the ant algorithm leaves information on the path for the sake of next routing during the search process. The information can be kept when more nodes are added. In the future, multi-robot path planning problem in dynamic obstacle environment will be studied. At that time the foregoing knowledge methodology will be used.

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